

2. List all the data structures (including the operations) that you will use to implement the Proposal Algorithm. (e.g. Stacks/Arrays/...). What is the complexity of your algorithm? Analyze in detail.
3. Consider the stable matching problem, where certain man-woman pairs are forbidden. Each man m ranks all the women, except the ones which are forbidden, and same with each woman. We can assume that the set F consists of the forbidden pairs, $F \subseteq M \times W$, where M is the set of n -men, and W is the set of n -women. Modify the definition of stable matchings to take into account the forbidden set F . Modify the above proposal algorithm to compute a set of stable matchings. Prove that your algorithm terminates and produces stable matchings.
4. Let $p(n) = a_d n^d + a_{d-1} n^{d-1} + \dots + a_1 n + a_0$, where $a_d > 0$, be a d -degree polynomial in n . Also a_0, \dots, a_d are constants. Let k be a positive integer. Prove the following:
 - (a) If $k \geq d$, then $p(n) = O(n^k)$.
 - (b) If $k \leq d$, then $p(n) = \Omega(n^k)$.
 - (c) If $k = d$, then $p(n) = \Theta(n^k)$.
 - (d) If $k > d$, then $p(n) = o(n^k)$.

(For this question please review the definitions of O, Θ, Ω, o , and prove each of the above statements by showing the appropriate constants c, c_1, c_2, n_0 ; See Chapter 3 in the book.)

5. Solve the recurrence relation

$$T(n) = T(xn) + T((1-x)n) + cn$$

in terms of x and n where x is a constant in the range $0 < x < 1$. Is the asymptotic complexity the same when $x = 0.5, 0.1$ and 0.001 ? What happens to the constant hidden in the $O()$ notation. (Hint: Please review the last part of Section 4.4 in the 3rd edition of the text-book or Section 4.2 from the 2nd edition.)

6. Show that the following recurrence evaluates to $O(n)$.

$$T(n) \leq \begin{cases} T(\lceil n/5 \rceil) + T(\lceil 7n/10 \rceil) + O(n) & \text{if } n > 140, \\ O(1) & \text{if } n \leq 140. \end{cases}$$

(This one is in the book!)

7. Analyze the recurrence relation

$$T(n) = \sqrt{n}T(\sqrt{n}) + n.$$

(This is called as the rootish-divide-and-conquer and for example plays an important role in parallel computing.)

8. Assume that you have a sorted array of size n containing real numbers. State the recurrence relation for time for performing binary search on this array. Analyze it and show that it evaluates to $O(\log n)$; Is it $\Theta(\log n)$ as well?
9. Suppose you need to choose between the following algorithms which solves the same problem:
- (a) Algorithm A solves the problem by dividing it into 5 subproblems of half of the size, recursively solves each of them, and combines the solution in linear time.
 - (b) Algorithm B solves the problem of size n by recursively solving two subproblems of size $n - 1$ and then combining the solutions in constant time.
 - (c) Algorithm C solves the problem of size n by dividing it into 9 subproblems of size $n/3$ each, recursively solving each of them, and then combining the solution in $O(n^2)$ time.

What are the running times of each of these algorithms and which one will you choose?

10. Solve the following recurrence $T(n) = 2T(n/2) + O(n \log n)$, where $T(n)$ is a constant for all values of $n \leq 4$.
11. (Bonus Problem) Can you devise a proposal algorithm that is unbiased? (Look up on Google and find an appropriate reference!)
12. (Bonus Problem) V. Pan has discovered a way to multiply two $70 * 70$ matrices using only 143640 multiplications. Ignoring the additions, what will the asymptotic complexity of Pan's algorithm for multiplying two $n * n$ matrices? Is it better than Strassen's? Justify your answer.